

Peer-to-Peer Data Sharing for Scientific Workflows on Amazon EC2

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Workflows in the Cloud

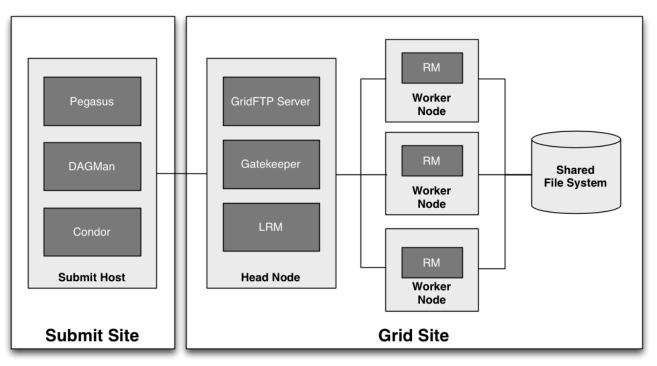
- Advantages
 - Provisioning (compute and storage)
 - Elasticity
 - Reproducibility
 - Appliances (e.g. Galaxy)
 - Control over environment (esp. for legacy)
- Disadvantages
 - Administration
 - Virtualization overhead
 - Resource limitations (not really infinite, no queuing)
 - Cost relative to alternatives (campus clusters, grid)
 - Cost/Performance tradeoffs





Deploying Workflows in the Cloud

- Could develop Workflow as a Service (PaaS or SaaS)
- Can deploy existing software on laaS clouds
- "Virtual Clusters"
- New tools: Nimbus Broker, cloudinit.d, Wrangler, Precip







Motivations for this Work

- Data-intensive workflows are limited by I/O performance
 - I/O is becoming the bottleneck rather than throughput
- Many workflows share data using files
 - Task A writes a file, task B reads it
 - File management is critical
- Write-once
 - Typically, files are only written once, never updated
 - Can replicate files without worrying about consistency
- Three ways to share files
 - 1. Use a shared storage system (POSIX or non-POSIX)
 - 2. Transfer files from submit host to workers and back
 - 3. Transfer files directly from one worker to the next





Previous Study on Data Sharing Options

- Goal
 - Better understand how storage systems affect performance
 - Compare storage costs on commercial clouds
- Deployed several different storage systems
 - Local, NFS, S3, PVFS2, GlusterFS (distribute and NUFA)
- Used three different workflow applications with different resource requirements
 - Montage (astronomy, data-intensive)
 - Broadband (seismology, memory-intensive)
 - Epigenome (bioinformatics, CPU-intensive)

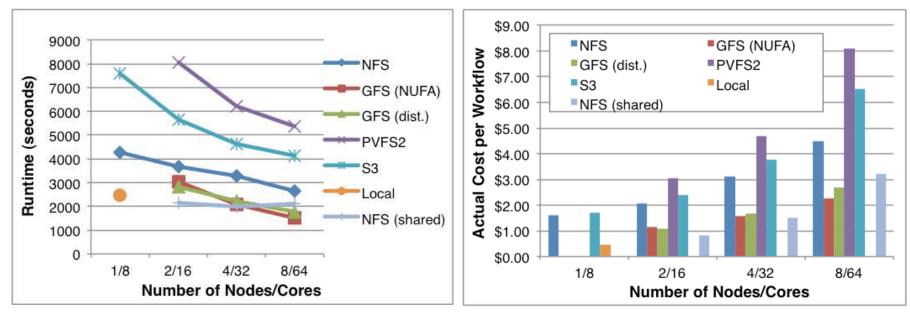
Compared performance and cost of different file system options

G. Juve, et al., "Data Sharing Options for Scientific Workflows on Amazon EC2", Supercomputing, 2010.





Results for Montage



Makespan

- PVFS didn't handle small files well
- S3 had too much overhead
- NFS did comparatively well
- GlusterFS came out on top

Cost

- NFS and S3 have extra costs
- Performance improvement does not offset increased cost





Approach

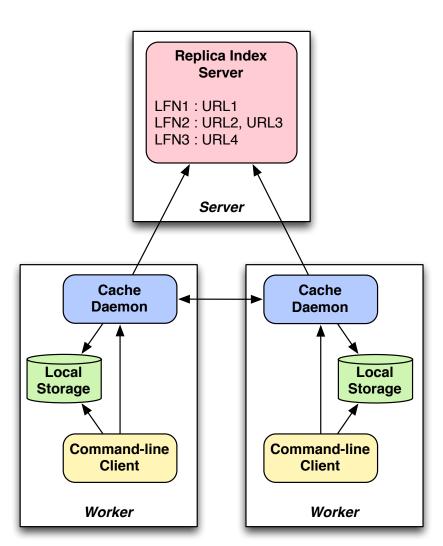
- Develop storage service to facilitate peer-to-peer transfers
 - Applies to environments other than clouds
- New files are written to the local disk
 - No network I/O for writes
- Files are replicated on-demand
 - Each time a task runs on a worker, all of its input files are replicated to that worker
- Files cached on each worker node
 - Enabled by write-once, no consistency issues
- Workflow tasks are wrapped by I/O operations
 - 1. Fetch input files
 - 2. Run task
 - 3. Register output files





System Design

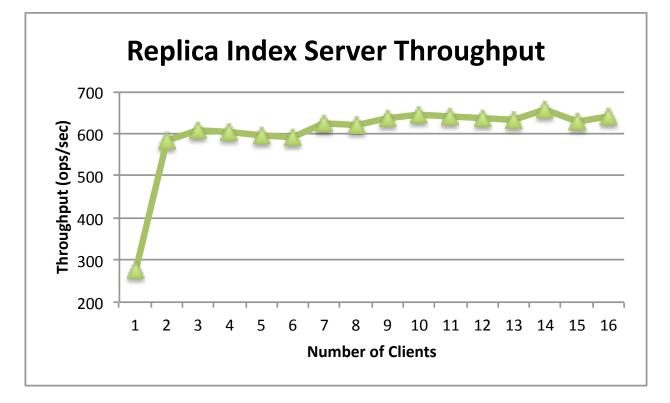
- Replica Index Server
 - Stores mappings of logical file names to URLs
- Cache Daemon
 - Manages local storage on each worker
 - Serves local replicas to peers
 - Retrieves remote replicas from peers
- Command-line Client
 - Get files from remote storage
 - Put files into local storage







Replica Index Server Throughput Benchmark



- Set up RIS on m1.xlarge, issued 1000 add operations each from 1-16 clients on m1.medium instances
- RIS achieved a peak throughput of ~650 ops/sec





Benchmarked vs. Observed RIS Throughput

Average requests per second observed for a 10-degree Montage workflow

Nodes / Cores	Entries in RIS	Workflow runtime (sec)	Average put requests/second			
2/8	63558	6699	9.5			
4 / 16	76688	4705	16.3			
8 / 32	N/A	3690	N/A			
16 / 64	87073	3704	23.5			

- Ran 10 degree workflow using 8-64 cores (m1.xlarge)
- Observed RIS throughput (10-25 ops/sec) is much less than benchmarked throughput (650 ops/sec)
- RIS should not be the bottleneck for workflows and resource pools of this size





Cache Daemon Benchmarks

	Implementation		0 MB	8 1 MB		10 M	В	100 MB
Put Latency (sec)	сору		0.0	07	0.00	09 ().35	4.36
	symlink		0.0	08	0.00	0.	008	0.008
	Implementation		0 MB	3 1 MB		10 ME	3	100 MB
Get Latency (sec)	сору		0.0	0.03		B1 0.1	178	3.951
	symlink		0.0	0.03		33 0. ⁻	146	1.841
	symlink+fsync		0.0	0.07		73 0.3	373	3.182
Get Bandwidth (MB/sec)		Implementation			1 MB	10 MB		00 MB
		сору			31.784	56.048		25.31
		symlink			30.571	68.734		54.329
		symlink+fsync			13.776	26.824		31.423

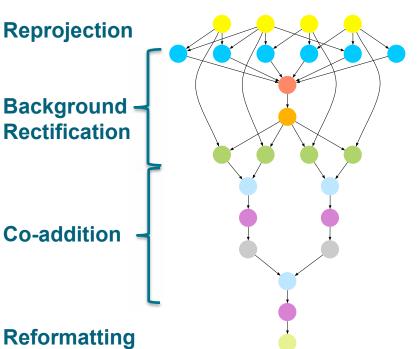
- Disk performance: ~38 MB/s write, ~109 MB/s read
- Network performance: ~89 MB/s
- Bottom line: Latency limits performance for smaller files

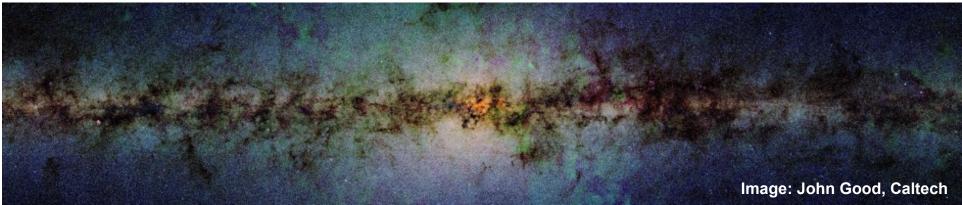




Workflow Performance Comparison

- Application: Montage
 - Creates science-grade astronomical image mosaics
- Test workflow
 - 10 degree square area
 - 19,320 tasks
 - 13 GB input, 88 GB output





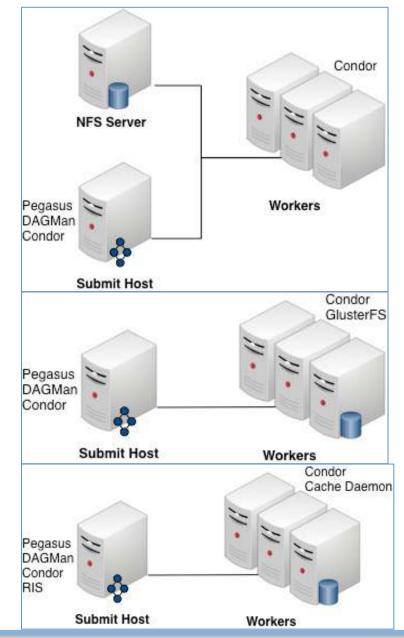


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Storage Systems

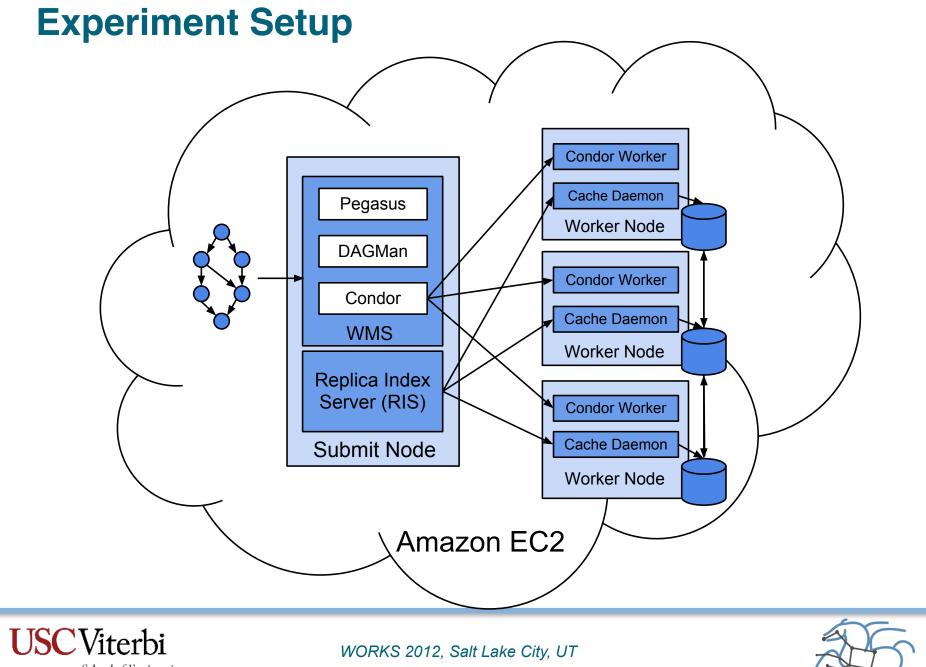
- NFS
 - Centralized file system
 - Used a dedicated m1.xlarge instance
- GlusterFS
 - Distributed file system
 - Used "distribute" mode
 - Each worker participates in the file system
- P2P
 - Our approach
 - RIS co-located with submit host





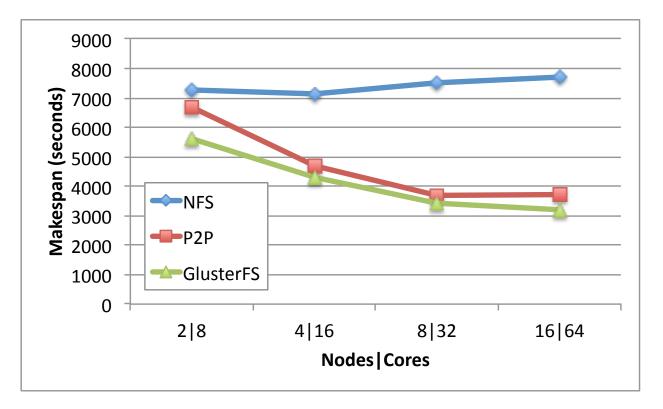
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Performance Comparison



- NFS performance is flat, as expected
- Performance flattens out due to workflow structure
- GlusterFS performs 13-16% better than P2P





Discussion

- Bottlenecks
 - Main problem with NFS
 - GlusterFS has no central server
 - P2P RIS is not a bottleneck based on benchmarks
- Latency
 - P2P query overhead harms small file performance
 - Not an issue for GlusterFS (just a hash to find the host)
- Load Balancing
 - P2P does not try to control data placement
 - GlusterFS distributes data more evenly
- Small reads
 - P2P always fetches the entire file
 - GlusterFS can fetch only the blocks required
 - Can overlap communication and computation



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Conclusion

- Our experiment did not work out as we hoped, but produced some valuable results
 - RIS server was not a bottleneck
 - Overheads were significant for small files
- We now have a better understanding of the problem
 - Partial reads may be important for some workflows
 - Locality and load balancing are important
 - Need to consider planning and scheduling data movement



Future Work

- Experiment with more workflows
- Compare with alternative data storage solutions
 - e.g. SRM, IRODS
- Study the I/O patterns of different workflows
 - e.g. partial reads
- Optimize the system, especially latencies
- Investigate techniques for planning data placement
- Make use of data-aware scheduling heuristics



